#### Localization

#### Sensor Networks

Prof. Dr. Kay Römer

# **Localization of Sensor Nodes**

- Estimation of node coordinates (x,y,z) in (Euclidean) coordinate system
- Why needed?
  - Interpretation of sensor data
  - Data fusion
  - Geo routing
  - ...
- Large design space
  - Intern vs. extern
  - always vs. sometimes
  - All nodes vs. some
  - Points vs. intervals

Уı  $(x_2, y_2)?$  $(x_4, y_4)?$ (x<sub>1</sub>, y<sub>1</sub>)?  $(x_3, y_3)?$ Χ

## **GPS**?

- GPS for sensor nodes?
- Example: U-blox ZOE-M8B SiP
  - Very small (4.5 x 4.5 x 1.0 mm)
  - Accuracy: 2.5m
  - Energy: 12-72 mW, 1.8 V
  - Startup time: 26 sec



 But: some nodes with GPS may act as anchor nodes for localization of others

- < 10% of nodes

 Localization task: Estimation of positions of remaining nodes



# **Localization Algorithms**

- Distributed localization algorithms with three phases
  - Measure relations between nodes
    - Distance (range)
    - Angle
    - ...
  - Initial positioning
    - Lateration
    - Angulation
    - ..
  - Refinement
    - Iterative refinement of initial positions

## **Node Relations**

- Mostly distance-based
  - Why not angles?
- Signals
  - Radio
  - (Ultra) sound
  - Light
- Indicators
  - Amplitude (signal strength)
  - Time of arrival
  - Frequency
  - Phase

# **Sound: Propagation Time**

- Measurement of propagation time of ultrasound from sender to receiver
  - Synchronization via radio
  - $s = t v_{sound}$
- Coding of sound signal
  - Robustness
  - Pseudo noise
- Receiver finds earliest occurrence of sound pattern
  - Echoes!



#### **Accuracy: Best Case**

- Accuracy ~ 1 cm
- Range ~ 10m
- Conditions
  - Free line of sight
  - Speaker and microphone facing each other



#### **Accuracy: Temperature**

- Sound propagation speed depends on air temperature
  - ~ 10% error per 50 K
  - Countermeasure: temperature sensor!



## **Accuracy: Orientation**

- Relative orientation of speaker and microphone has big impact
  - Countermeasure: Scattering reflectors, multiple speakers / micros



### **Accuracy: Obstacles**

- Light obstacles (e.g., card box)
  - Sound travels through / around obstacle
  - Small offset
- Massive obstacles (e.g., mattrass)
  - Multipath
  - Big offset
- Countermeasure: Test for free LOS (e.g., using light)



## BeepBeep

- Idea: both devices beep with a constant known delay and they record their own as well as the other's beep
- Time is measured by sample counting since continuous sampling is utilized
- No time sync needed!



# **Received Signal Strength**

 RSS is function of distance and channel between sender and receiver

- Radio outputs P<sub>r</sub> as RSS indicator
  - Received Signal Strength Indicator [dBm]  $\sim RSSI = 10 \log_{10}(P_r / 1mW)$

$$RSSI = 10\log_{10} P_s - 10n\log_{10} d$$

 $0 \, dBm =$ 

 $-10 \, dBm =$ 

 $-20 \, dBm =$ 

-30 dBm = 0.001 mW

1 mW

0.1 mW

0.01 mW

n=2

RSS

## **Transmitted Signal Strength**

- Indirect measurement of transmitted signal strength at reference distance d<sub>0</sub>
  - Dependent on antenna gain

$$P_r(d_0) \approx P_s \frac{1}{d_0^n} \qquad P_s \approx P_r(d_0) d_0^n$$

 $\left| \text{RSSI} = 10 \log_{10} P_r(d_0) + 10n \log_{10} d_0 / d \right|$ 

- P<sub>r</sub>(d<sub>0</sub>) typically given in the data sheet
  - cc2420:  $10\log_{10}P_r(2m) = -46dbm$  at transmit power 0dBm

## **Accuracy: Best Case**

- Accuracy of 2-3m possible
- Range 10-100m
- Conditions
  - Free space, no walls, buildings, objects etc.
  - Nodes at equal height above ground
  - Large distance from ground
  - Aligned antennas



## **Accuracy: Orientation**

- Relative orientation of antennas has big impact
- Distance above ground has big impact





#### **Accuracy: Indoor** In buildings strong multipath effects - Distance and RSSI almost uncorrelated Can be exploited for detecting moving obstacles and changing environments Distance Vs Average RSSI at the maximum power level 6.17ft -105.65ft -15

-20



#### Why is RF Ranging Challenging?





Two nodes 30m apart at varying heights



- At low angles, reflection has a phase shift of PI, distance difference between LOS and ground reflected signal is minimal: destructive interference
- Any additional even weak multipath can affect amplitude and phase significantly

# Connectivity

- RSS often not useful
- Binary distance measure?
  - d=1: Within communication range
  - $d=\infty$ : Outside communication range
- Distance measured by number of hops



## **Radio Interferometry**

- A and B transmit carrier signal at slightly different frequencies
  - Few 100 Hz
- Phase of interfered signal depends on receiver location
- Measurement of  $d_{ABCD}$ 
  - Measure phase offset at multiple frequencies



- Solve equation system





## **Radio Interferometry**

- Very accurate outdoors
  - Few cm
- Range up to 160 m
- Requires
  - Accurate time sync of C and D: 1 us
  - Accurate tuning of transmitter frequencies
- "Strange" distance measure
  - Solve complex optimization problem
  - Find node locations that minimize difference from measured values

# **Radio Interferometric AOA**

- Motivation: Measure AOA with motes without any additional hardware
- Group 3 nodes to form an anchor array
  - Orthogonal antennas to minimize parasitic effects
  - Array uses radio interferometry to estimate bearing to target node





# **Initial Positions**

#### Input

- Anchors with known locations
- Relations among nodes
- Output
  - Position of non-anchor nodes
- Challenges
  - Inaccurate relations
  - Relations only among neighbors
  - Sparse anchors
  - Arrangement of anchor nodes





## Centroid

 Position node at centroid of all anchors within communication range

$$x = \frac{1}{n} \sum_{i=1}^{n} x_i$$
  $y = \frac{1}{n} \sum_{i=1}^{n} y_i$ 



## **Bounding Box**

- Position node at center point of intersection of bounding boxes around anchors i=1..N with position (x<sub>i</sub>, y<sub>i</sub>) and distance d<sub>i</sub> x = <sup>1</sup>/<sub>2</sub>(min<sub>i</sub>(x<sub>i</sub> + d<sub>i</sub>) + max<sub>i</sub>(x<sub>i</sub> - d<sub>i</sub>))
  - $y = \frac{1}{2}(\min_{i}(y_{i} + d_{i}) + \max_{i}(y_{i} d_{i}))$
- Variant: Weighted sum of corners j

$$W(j) = \frac{1}{\sum_{i=1..N} (D_{ij} - d_i)^2}$$





## Lateration

Position node at intersection of circles around anchors

$$(x - x_{1})^{2} + (y - y_{1})^{2} = d_{1}^{2}$$

$$(x - x_{2})^{2} + (y - y_{2})^{2} = d_{2}^{2}$$

$$(x - x_{3})^{2} + (y - y_{3})^{2} = d_{3}^{2}$$

$$2(x_{1} - x_{2})x + 2(y_{1} - y_{2})y = d_{2}^{2} - d_{1}^{2} + (x_{1}^{2} + y_{1}^{2}) - (x_{2}^{2} + y_{2}^{2})$$

$$2(x_{1} - x_{3})x + 2(y_{1} - y_{3})y = d_{3}^{2} - d_{1}^{2} + (x_{1}^{2} + y_{1}^{2}) - (x_{3}^{2} + y_{3}^{2})$$

$$\begin{pmatrix} a_{11} & a_{12} \\ a_{21} & a_{22} \end{pmatrix} \begin{pmatrix} x \\ y \end{pmatrix} = \begin{pmatrix} b_{1} \\ b_{2} \end{pmatrix}$$

$$\boxed{A\begin{pmatrix} x \\ y \end{pmatrix} = b}$$

$$27$$

## Lateration

- Problem: inaccurate distances
  - No common intersection point
  - Equation system does not have a solution!
- Approach: formulate optimization problem



## **Multilateration**

- Same approach, but n instead of 3 anchors
  - A has n-1 rows
  - A<sup>T</sup>A still 2x2 matrix
  - A<sup>T</sup>b still vector of length 2
- Similar for 3D

$$\left[ \left( A^T \cdot A \right) \begin{pmatrix} x \\ y \end{pmatrix} = A^T b \right]$$

# **Intersecting Circle Pairs**

- Pairwise intersection of circles around anchors
  - If no intersection use midpoint of shortest line connecting the circles
- Remove outlier intersections
  - Not contained in most (< N-2) circles around anchors
  - Sum of distances to other intersections larger than median
- Compute centroid of remaining intersections

#### **Bounding Box vs. Lateration**

- Bounding Box better as inaccuracy of ranging grows
- Anchor arrangement important



Lateration



**Bounding Box** 



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# **Multi-Hop Localization**

- Up to now: every node needs 3 anchor neighbors
  - Or even more for high accuracy
- Approaches
  - Dense anchors
  - Multi-Hop relations
    - Distance to anchors that are not neighbors
  - Recursion
    - Positioned nodes as additional anchors



# **Multi-Hop Relations**

 Compute shortest paths (Euclidean) to anchors

 $d_3$ 

 $d_2$ 

**d**₁

- $D = d_1 + d_2 + d_3$
- Never smaller than true distance
- Zigzag paths lead to error
- Implementation
  - Anchors flood network
  - Nodes compute and rebroadcast shortest distance to anchor

# **Multi-Hop Relations**

- Estimate average hop distance
   L = X / (1+1)
- Compute shortest paths (Hops) to anchors
  - D = (1+1+1) x L
  - No bounds
  - Zigzag
- Implementation
  - Anchors flood network
  - Anchors compute hop-distance among each other and flood again

#### Recursion

- Use newly positioned nodes as additional anchors
  - Errors accumulate
  - May get stuck



#### **Multi-Hop: Accuracy**



#### **Multi-Hop: Coverage**



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## **Iterative Improvement**

- Nodes have initial positions
- Iterative improvement
  - Each node recomputes position using all neighbors as anchors
  - Iterate until positions converge to a fixed point
- May not converge

#### **Iterative Improvement**



# **Other Approaches**

- Centralized algorithms
  - Mass-spring models (cf. graph layout)
  - Convex optimization
  - Genetic optimization
- Probabilistic algorithms
- Optimizing anchor locations
- Anchor-free approaches
- Support for mobility
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## References

- Slides contain material by following authors
  - Lewis Girod UCLA
  - Dimitrios Lymberopoulos Yale
  - Branislav Kusy Vanderbilt
  - Alec Woo Berkeley
  - Nirupama Bulusu Portland State
  - Koen Langendoen TU Delft
  - Akos Ledeczi Vanderbilt