Localization

Sensor Networks

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Localization of Sensor Nodes

- **Estimation of node coordinates** (x, y, z) **in** (Euclidean) coordinate system
- Why needed?
	- Interpretation of sensor data
	- Data fusion
	- Geo routing
	- ...

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- **Large design space**
	- Intern vs. extern
	- always vs. sometimes
	- All nodes vs. some
	- Points vs. intervals

1 (x_2, y_2) ? $\mathbf{2}^{\setminus}$ 3 4 (x_4, y_4) ? (x_1, y_1) ? $\frac{y_1}{(x_3, y_3)}$? y. x

GPS?

- GPS for sensor nodes?
- Example: U-blox ZOE-M8B SiP
	- Very small $(4.5 \times 4.5 \times 1.0 \text{ mm})$
	- Accuracy: 2.5m
	- Energy: 12-72 mW, 1.8 V
	- Startup time: 26 sec

 But: some nodes with GPS may act as anchor nodes for localization of others

 $- < 10\%$ of nodes

ELOCAlization task: Estimation of positions of remaining nodes

Localization Algorithms

- Distributed localization algorithms with three phases
	- Measure relations between nodes
		- Distance (range)
		- Angle
		- ...
	- Initial positioning
		- Lateration
		- Angulation
		- ...

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- Refinement
	- Iterative refinement of initial positions

Node Relations

- Mostly distance-based
	- Why not angles?
- **Signals**
	- Radio
	- (Ultra) sound
	- Light
- **Indicators**
	- Amplitude (signal strength)
	- Time of arrival
	- Frequency
	- Phase

Sound: Propagation Time

- Measurement of propagation time of ultrasound from sender to receiver
	- Synchronization via radio
	- $s = t v_{sound}$
- **Coding of sound signal**
	- **Robustness**
	- Pseudo noise
- **Receiver finds earliest occurrence of sound pattern**
	- Echoes!

Accuracy: Best Case

- Accuracy ~ 1 cm
- Range ~ 10m
- **Conditions**
	- Free line of sight
	- Speaker and microphone facing each other

Accuracy: Temperature

- **Sound propagation speed depends on** air temperature
	- ~ 10% error per 50 K
	- Countermeasure: temperature sensor!

Accuracy: Orientation

- Relative orientation of speaker and microphone has big impact
	- Countermeasure: Scattering reflectors, multiple speakers / micros

Accuracy: Obstacles

- Light obstacles (e.g., card box)
	- Sound travels through / around obstacle
	- Small offset
- Massive obstacles (e.g., mattrass)
	- Multipath
	- Big offset
- Countermeasure: Test for free LOS (e.g., using light)

BeepBeep

- Idea: both devices beep with a constant known delay and they record their own as well as the other's beep
- **Time is measured by sample counting since continuous** sampling is utilized
- No time sync needed!

Received Signal Strength

 RSS is function of distance and channel between sender and receiver

$$
P_r(d) \approx P_s \frac{1}{d^n} \cdot \frac{2 \cdot 2 \cdot 6!}{e^{n-2} \cdot 6!}
$$

- Radio outputs P_r as RSS indicator
	- Received Signal Strength Indicator [dBm] $\cos RSSI = 10 \log_{10} (P_r /1mW)$

-20 dBm = 0.01 mW -30 dBm = 0.001 mW -40 dBm = 0.0001 mW

 $\overline{0}$ dB 10 d $\overline{1}$

Transmitted Signal Strength

- **Indirect measurement of transmitted** signal strength at reference distance d_0
	- Dependent on antenna gain

$$
P_r(d_0) \approx P_s \frac{1}{d_0^n} \qquad P_s \approx P_r(d_0) d_0^n
$$

 $RSSI = 10 \log_{10} P_r(d_0) + 10 n \log_{10} d_0/d$

- \blacksquare P_r(d₀) typically given in the data sheet
	- cc2420: $10\log_{10}P_r(2m) = -46d$ bm at transmit power 0dBm

Accuracy: Best Case

- Accuracy of 2-3m possible
- Range 10-100m
- **Conditions**
	- Free space, no walls, buildings, objects etc.
	- Nodes at equal height above ground
	- Large distance from ground
	- Aligned antennas

Accuracy: Orientation

- **Relative orientation of** antennas has big impact
- **Distance above ground** has big impact

Accuracy: Indoor

In buildings strong multipath effects

- Distance and RSSI almose uncorrelated

Why is RF Ranging Challenging?

- At low angles, reflection has a phase shift of PI, distance difference between LOS and ground reflected signal is minimal: destructive interference
- Any additional even weak multipath can affect amplitude Two nodes 30m apart at varying heights **and phase significantly**

Connectivity

- RSS often not useful
- **Binary distance measure?**
	- d=1: Within communication range
	- d=∞: Outside communication range
- Distance measured by number of hops

Radio Interferometry

- **A and B transmit carrier** signal at slightly different frequencies
	- Few 100 Hz
- **Phase of interfered signal** depends on receiver location
- Measurement of d_{ABCD}
	- Measure phase offset at multiple frequencies

Radio Interferometry

- Very accurate outdoors
	- Few cm
- Range up to 160 m
- **Requires**
	- Accurate time sync of C and D: 1 us
	- Accurate tuning of transmitter frequencies
- **"** "Strange" distance measure
	- Solve complex optimization problem
	- Find node locations that minimize difference from measured values

Radio Interferometric AOA

- Motivation: Measure AOA with motes without any additional hardware
- Group 3 nodes to form an anchor array
	- Orthogonal antennas to minimize parasitic effects
	- Array uses radio interferometry to estimate bearing to target node

Initial Positions

Input

- Anchors with known locations
- Relations among nodes
- Output
	- Position of non-anchor nodes
- Challenges
	- Inaccurate relations
	- Relations only among neighbors
	- Sparse anchors
	- Arrangement of anchor nodes

Centroid

 Position node at centroid of all anchors within communication range

$$
x = \frac{1}{n} \sum_{i=1}^{n} x_i \qquad y = \frac{1}{n} \sum_{i=1}^{n} y_i
$$

Bounding Box

- **Position node at center point of intersection** of bounding boxes around anchors i=1..N with position ($\mathsf{x}_{\mathsf{i}},\, \mathsf{y}_{\mathsf{i}}$) and distance d_{i} $\frac{1}{2}$ (min_i(y_i + d_i) + max_i(y_i - d_i)) $\frac{1}{2}$ (min_i(x_i + d_i) + max_i(x_i - d_i)) 1 1 $y = \frac{1}{2}(\min_i(y_i + d_i) + \max_i(y_i - d_i))$ $x = \frac{1}{2}(\min_i(x_i + d_i) + \max_i(x_i - d_i))$
- Variant: Weighted sum of corners j

$$
W(j) = \frac{1}{\sum_{i=1..N} (D_{ij} - d_i)^2}
$$

$$
x = \frac{\sum W(j) \cdot ex_j}{\sum W(j)} \qquad y = \frac{\sum W(j) \cdot ey_j}{\sum W(j)}
$$

Lateration

Position node at intersection of circles around anchors

Lateration

- **Problem: inaccurate distances**
	- No common intersection point
	- Equation system does not have a solution!
- Approach: formulate optimization problem

Multilateration

- Same approach, but n instead of 3 anchors
	- A has n-1 rows
	- ATA still 2x2 matrix
	- ATb still vector of length 2
- Similar for 3D

$$
(AT \cdot A) \begin{pmatrix} x \\ y \end{pmatrix} = AT b
$$

Intersecting Circle Pairs

- **Pairwise intersection of circles around anchors**
	- If no intersection use midpoint of shortest line connecting the circles

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3

2

- Remove outlier intersections
	- Not contained in most (< N-2) circles around anchors
	- Sum of distances to other intersections larger than median
- Compute centroid of remaining intersections

Bounding Box vs. Lateration

Bounding Box better as inaccuracy of ranging grows

100

80

60

40

20

0

0

 0.05

 0.1

Range error stddev

Position error

Anchor arrangement important

Lateration

 0.15

BBox

 0.2 31

 0.25

Lateration

Bounding Box

Multi-Hop Localization

- Up to now: every node needs 3 anchor neighbors
	- Or even more for high accuracy
- Approaches
	- Dense anchors
	- Multi-Hop relations
		- Distance to anchors that are not neighbors
	- Recursion
		- Positioned nodes as additional anchors

Multi-Hop Relations

- Compute shortest paths (Euclidean) to anchors
	- $D = d_1 + d_2 + d_3$
	- Never smaller than true distance
	- Zigzag paths lead to error
- **Implementation**
	- Anchors flood network
	- Nodes compute and rebroadcast shortest distance to anchor

 O_4

 d_2

 d_3

Multi-Hop Relations

- **Estimate average hop distance** $-L = X / (1+1)$
- Compute shortest paths (Hops) to anchors
	- $-D = (1+1+1) \times L$
	- No bounds
	- Zigzag
- **Implementation**
	- Anchors flood network
	- Anchors compute hop-distance among each other and flood again

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Recursion

- Use newly positioned nodes as additional anchors
	- Errors accumulate
	- May get stuck

Multi-Hop: Accuracy

Multi-Hop: Coverage

Iterative Improvement

- **Nodes have initial positions**
- **Interative improvement**
	- Each node recomputes position using all neighbors as anchors
	- Iterate until positions converge to a fixed point
- May not converge

Iterative Improvement

Other Approaches

- Centralized algorithms
	- Mass-spring models (cf. graph layout)
	- Convex optimization
	- Genetic optimization
- **Probabilistic algorithms**
- **Optimizing anchor locations**
- **Anchor-free approaches**
- **Support for mobility**
- ...

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References

- **Slides contain material by following** authors
	- Lewis Girod UCLA
	- Dimitrios Lymberopoulos Yale
	- Branislav Kusy Vanderbilt
	- Alec Woo Berkeley
	- Nirupama Bulusu Portland State
	- Koen Langendoen TU Delft
	- Akos Ledeczi Vanderbilt